# Data Management and File Organization

Sequential Files

Part one: Pile files



## **Topics**

- Definitions:
  - Entity
  - Attribute
  - Record
- Files, Sequential Files
- Sequential File Operations and Timings



#### **Review: Motivation**

- Large data cannot be stored in the main memory of computers because:
  - The capacity of the main memory (RAM) is small
  - The data in the main memory is lost when we turn-off the computer
- Solution:
  - Using files stored on hard disks



### Questions:

- What is a file?
- How do we put data in a file?
- Is there any structure in the data stored in a file?
- How can we make file access fast?

We will find answers to these questions

### **Basic Definitions**



- Entity: Anything that we store some data about
- Attribute: Any property that helps to identify an entity.
   Ex. Name, height, age
- Data: Any value given to an attribute, Ex. 'Ali' for attribute Name
- Record: A group of attribute values which uniquely identifies an entity.

Ex. For entity 'student', the record is Student <Student ID, name, surname, major, year started, address, phone>

<200511777, Ali, Yildiz, Computer Eng., 2005, Ankara, 1234567 >



#### **Basic Definitions**

• File: A set of related records.

Ex. Student file (Set of student records)
Hospital file (Set of patient records)

**Exceptions**: Text files and binary files



## Sequential Files

 Definition: A sequential file is a file which is read from beginning to end.

- Types:
  - Unsorted sequential files (Pile files): A set of records with no order
  - Sorted sequential files: The records are sorted in the order of an attribute



#### File Header

Header is part of the file which includes data such as:

- Record size
- Block size
- Number and type of indexes
- Address of the last block

Header is read into memory when the file is opened

# Sequential File Operations and Timings



Fetch one record

 $T_{\mathsf{F}}$ 

Fetch next record

 $\mathsf{T}_{\mathsf{N}}$ 

Insert a record

 $\mathsf{T}_\mathsf{L}$ 

Update a record

 $\mathsf{T}_{\mathsf{U}}$ 

Delete a record

 $\mathsf{T}_\mathsf{D}$ 

Exhaustive reading of the file

 $\mathsf{T}_\mathsf{X}$ 

Re-Organize a file

 $\mathsf{T}_\mathsf{Y}$ 



#### Fetch One Record

- Find and read a record given an attribute value.
  - Example. Find student record with Student ID=200612345
- In a pile file on average half of the blocks are read
- T<sub>F</sub> = s + r + b/2\*ebt
   b: number of blocks in the file



# Example

Find T<sub>F</sub> given:

- Total number of records = 100,000
- Blocking factor (Bfr) = 6
- Number of blocks (b) = 100,000/6 = 16667
- ebt = 0.84 msec
- s=16msec
- r=8.3msec



#### Fetch Next Record

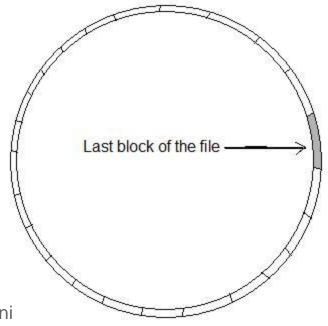
 Find and read the next record in order of an attribute value.

If the file is not sorted T<sub>N</sub> = T<sub>F</sub> (Each fetch is independent)

#### Insert a New Record



- Insert is always done at the last block
  - Time to read the last block = s + r + btt
  - Time to write back the block = 2r
  - $T_1 = s + r + btt + 2r$



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## Update a Record

- To update, first the block is read, then the record is updated and the block is written back
- Time to read the block =  $s + r + b/2*ebt = T_F$
- Time to write back the block = 2r
- $T_{U} = T_{F} + 2r$



#### Delete a Record

- To delete a record, we mark it as deleted
- First read the block (s + r + b/2\*ebt )
- Update the mark and write the block (2r)
- $T_D = T_F + 2r$

Mark	Record
1	Rec1
0	Rec2
0	Rec3
1	Rec4



# Exhaustive Reading of a File

Case 1: Read records without any attribute order
 T<sub>X</sub> (Beginning to End) = b\*ebt + s + r

#### Example:

```
b = 16667

ebt = 0.84 \text{ msec}

s=16 \text{ msec}
```

r=8.3 msec



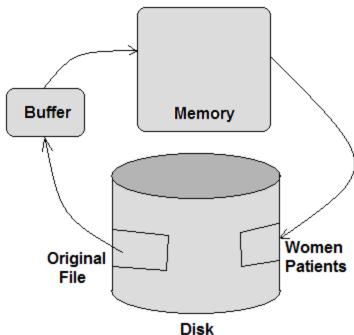
## Example

 Assume a hospital file records include illness, and duration of stay in hospital too. To find the average duration of stay in hospital for each illness, an exhaustive reading of the file without any attribute order is used



# Example

 Write records of the hospital file which show a female patient in a new file. An exhaustive reading without any attribute order is used here too.





## Exhaustive Reading of a File

- Case 2: with the order of an attribute value
- Tx (order of an attribute) =
   Number of records x T<sub>N</sub>



# Example

- Find the time needed to read the hospital file with 100,000 records.
- Assume
  - Bfr = 10
  - ebt = 0.84 msec
  - s=16 msec
  - r=8.3 msec



#### Solution

- Case 1: Without any attribute order
  - Tx = b\*ebt + s + r
  - b = 100,000 /10 = 10,000 (number of blocks, Bfr is 10)
  - Tx = 10,000 \* 0.84 + 16 + 8.3 = 8.424 s



#### Solution

- Case 2: with the order of an attribute
  - $Tx = 100,000 * T_N$
  - Assume  $T_N = 4224$  ms
  - About 4.89 days





# Re-Organizing a File

- In delete operation records are not physically deleted from the file.
- Operations in a file with many marked records are slow
- In re-organizing a file, marked records are deleted physically by copying active records to a new file.
- Old file is deleted



## Re-organizing files

- Time to read the file= s + r + b\*ebt
- Time to write active records=s+r+(n/Bfr)\*ebt
   (n is the number of active records)
   T<sub>Y</sub> = s+r+b\*ebt + s+r+(n/Bfr)\*ebt
- If two disks are available then
  - Read from disk one and write to disk two at the same time
  - $-T_Y = s+r+b*ebt$





- Assume a pile file has 100,000 records in it. Bfr=5, and 25% of the records are marked as deleted. Find T<sub>F</sub> for this file.
- Now assume the file has been reorganized. Find T<sub>F</sub> again and compare with your answer before re-organizing the file.
- Use: s=16msec, r=8.3msec , ebt = 0.84msec



## Summary



- File is a set of related records
- Sequential files are read from beginning to end
- File I/O operation timings depend on:
  - File size
  - Blocking factor
  - Order of reading records

## Questions?

